

## **METHOD FOR SECURELY MANAGING INFORMATION IN DATABASE**

### **BACKGROUND OF THE INVENTION**

#### **1. Field of the Invention**

The present invention relates to a method for securely managing information in a database, and more particularly, to a method for  
5 storing and securely managing important information for a user in a database.

#### **2. Description of the Related Art**

FIG. 1 is a diagram of a related art communication system, for example, an Internet system, for Digital Right Management (DRM).  
10 User computers 10a...10d, a content providing server 141, a key-data management server 142, and a payment system 143 are connected to one another through a communication network 12. The content providing server 141 provides content to the user computers 10a...10d and receives charges related to DRM through the payment system  
15 143. The key-data management server 142 manages key-data for encrypting and decrypting content files, which the content providing server 141 provides to the user computers 10a...10d.

Referring to FIG. 2, a related art content file includes a header 31 and data 32. The header 31 includes encryption information EI, user key-data KU, and DRM information DRM. The encryption information EI contains parameters used for encrypting the data 32, and the user key-data KU is used for encrypting and decrypting the data 32 of the content file. The DRM information indicates the number of use times permitted to a user.

However, the related art system and method has various problems and disadvantages. For example, but not by way of limitation, according to a simple method of managing information with such a related art file structure, the header 31 can be decrypted by a hacker even if it has been abstrusely encrypted. Accordingly, the hacker can change important information for a user such as DRM information. Therefore, in addition to encryption of information, it is necessary to efficiently determine whether the encrypted information has been accessed and/or altered without authorization (i.e., hacked) to securely manage information in a database.

### SUMMARY OF THE INVENTION

To solve at least the above-described problems, it is an object of the present invention to provide a method for securely managing information in a database by encrypting the information and efficiently determining whether the encrypted information has been hacked.

To achieve the above object of the invention, there is provided a method for storing important information for a user in a database and securely managing the stored important information for a user is provided. The method includes an encrypting step, a storing step, and  
5 a checking step. In the encrypting step, synchronizing information, which changes whenever at least one piece of the important information for a user is updated, is stored in the database together with the important information for a user and the synchronizing information is encrypted. In the storing step, the encrypted  
10 synchronizing information is distributively stored in predetermined places. In the checking step, the synchronizing information stored in the predetermined places is combined and decrypted, and it is checked whether the combined and decrypted synchronizing information is the same as the synchronizing information stored in the database.

15 According to the present invention, the synchronizing information stored in the predetermined places is combined and decrypted and whether it is the same as the synchronizing information stored in the database is checked, to efficiently determine whether the important information for a user has been hacked. Therefore, the  
20 present invention can manage the information in the database more securely than in the related art case of depending solely on encryption. In addition, since the present invention distributively stores the

encrypted synchronizing information in predetermined places, the risk of exposing the encrypted synchronizing information is minimized.

Additionally, to achieve at least the aforementioned objects and overcome at least the aforementioned problems and disadvantages of the related art, a method for storing and securely managing important  
5 information for a user in a database is provided, comprising (a) storing synchronizing information, which changes when at least one piece of the important information for a user is updated, in the database together with the important information for a user, and encrypting the  
10 synchronizing information, and (b) distributively storing the encrypted synchronizing information in a plurality of predetermined places. The method also comprises (c) combining and decrypting the synchronizing information stored in the predetermined places and determining whether the combined synchronizing information and the decrypted  
15 synchronizing information is substantially identical to the synchronizing information stored in the database.

Further, a content file is provided that comprises a header portion having key-data for synchronizing information and synchronizing information distributively stored in a plurality of  
20 predetermined places of a hard disc, and a data portion.

Also, a content file is provided that comprises a header portion having key-data for digital right management (DRM) information and



# DETAILED DESCRIPTION OF THE INVENTION

FIG. 3 illustrates a content file according to a preferred embodiment of the present invention that includes a header 41 and data 42. The header 41 includes encryption information EI, user key-data KU, key-data KS for synchronizing information, key-data KD for Digital Right Management (DRM) information, synchronizing information SI, and DRM information, which is important information for a user. In the present invention, "important information" is defined to include, but is not limited to, DRM information that indicates a number of use times permitted to a user, so as to enable effective performance of DRM.

In the header 41, the encryption information EI contains parameters used for encrypting the data 42 of the content file, and the user key-data KU is used for encrypting and decrypting the data 42 of the content file. Also, the key-data KS for synchronizing information is used for encrypting and decrypting the synchronizing information SI, and the key data KD for DRM information is used for encrypting and decrypting the DRM information. The synchronizing information SI is continuously updated to determine whether the important DRM information has been hacked. In a preferred embodiment of the present invention, the synchronizing information SI can be embedded in the DRM information. The important DRM information indicates the number of use times permitted to a user, but is not limited thereto.

FIG. 4 illustrates the encrypted synchronizing information SI, which is distributively stored in predetermined places. For example, but not by way of limitation, if an information managing method according to the present invention is applied only to the user computer 10a of FIG. 1 (here, the important information will not be the DRM information DRM of FIG. 3), the encrypted synchronizing information SI is distributed and stored in predetermined places of a hard disc drive. In this case, when the information managing method is applied to a communication network system as shown in FIG. 1, the encrypted synchronizing information SI may be distributed and stored in the key-data management server 142. Since the encrypted synchronizing information SI is distributed and stored in predetermined places, the risk of exposing the encrypted synchronizing information SI is minimized. According to this principle, the synchronizing information SI and the key-data KS for synchronizing information of FIG. 2 are distributed and stored in predetermined places according to a method shown in FIG. 5. The important DRM information DRM and the key-data KD for DRM information can be distributed and stored in predetermined places according to a method shown in FIG. 6.

FIG. 5 shows a first preferred embodiment of a program performed by the user computer 10a of FIG. 1 to securely manage the important DRM information DRM of the content file of FIG. 3. In the program of FIG. 5, the important DRM information DRM is not

encrypted, so the DRM information key-data KD of FIG. 3 is not necessary.

Key-data for synchronizing information distributively stored in predetermined places is combined and decrypted in step S501. Then,  
 5 synchronizing information distributively stored in predetermined places is combined and decrypted in step S502. Here, the synchronizing information is decrypted based on the key-data decrypted in step S501.

Next, encrypted synchronizing information SI of FIG. 3 stored in a database, for example, a hard disk within the user computer 10a of  
 10 FIG. 1, is decrypted in step S503. Next, in step S504, it is determined whether the synchronizing information decrypted in step S502 is the same as the synchronizing information decrypted in step S503. If the synchronizing information decrypted in step S502 is not the same as the synchronizing information decrypted in step S503, it can be  
 15 concluded that a content file including the DRM information DRM of FIG. 3 has been hacked. Accordingly, hacking information is displayed in step S506 and is transmitted to a server, for example, the content providing server 141 or the key-data management server 142 of FIG. 1, in step S507.

20 If the synchronizing information decrypted in step S502 is the same as the synchronizing information decrypted in step S503, then in step S504, update of the important DRM information DRM is monitored in step S508. The update of the DRM information DRM is related to, for



example, the number of uses of the data 42 of the content file of FIG.

3.

If the DRM information DRM has been updated, new  
synchronizing information and important DRM information DRM is  
5 stored in the database in step S509. Next, the synchronizing  
information SI stored in the database is encrypted in step S510. Next,  
key-data KS for the encrypted synchronizing information SI is  
encrypted in step S511. Sequentially, the encrypted synchronizing  
information and key-data is distributed and stored in predetermined  
10 places in step S512. Steps S508 through 512 are repeated until a  
termination signal is input in step S513.

FIG. 6 shows a second preferred embodiment of a program  
performed by the user computer 10a of FIG. 1 to securely manage the  
important DRM information DRM of the content file of FIG. 3. In the  
15 program of FIG. 6, the important DRM information DRM is encrypted,  
so the DRM information key-data KD of FIG. 3 is used.

Key-data for synchronizing information and key-data for DRM  
information defined as important information for a user, which are  
distributively stored in predetermined places, are combined and  
20 decrypted in step S601. Next, synchronizing information and the  
important DRM information, which are distributively stored in  
predetermined places, are combined and decrypted in step S602. In  
step S602, the synchronizing information and the DRM information are

decrypted based on the key-data decrypted in step S601. Next,  
synchronizing information SI and important DRM information DRM of  
FIG. 3, which is encrypted and stored in a database, for example, a  
hard disk of the user computer 10a of FIG. 1, is decrypted in step  
5 S603.

Sequentially, in step S604 it is determined whether the  
synchronizing information decrypted in step S602 is the same as the  
synchronizing information decrypted in step S603 and whether the  
important DRM information DRM decrypted in step S602 is the same  
10 as the important DRM information DRM decrypted in step S603. If the  
information is not the same in step S604, it can be concluded that a  
content file including the DRM information DRM has been hacked.  
Accordingly, hacking information is displayed in step S606 and is  
transmitted to a server, for example, the content providing server 141  
15 or the key-data management server 142 of FIG. 1 in step S607.

If the information is the same in step S604, update of the  
important DRM information DRM is monitored in step S608. The  
update of the DRM information DRM is related to, for example, the  
number of uses of the data 42 of the content file of FIG. 3. If the DRM  
20 information DRM has been updated, new synchronizing information  
and important DRM information DRM is stored in the database in step  
S609. Next, the synchronizing information SI and the important DRM  
information DRM of FIG. 3 stored in the database are encrypted in step

